Design and Performance analysis of a 3GPP LTE/LTE-Advance turbo decoder using software reference models

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Abstract— This paper presents the design and development of an efficient VLSI architecture for 3GPP advanced Turbo decoder by utilizing the convolutional interleaver. The high-throughput 3GPP Advance Turbo code requires turbo decoder architecture. Interleaver is known to be the main obstacle to the decoder implementation and introduces latency, due to the collisions it introduces in accesses to memory. In this paper, we propose a low-complexity soft Input Soft Output (SISO) turbo decoder for memory architecture to enable the Turbo decoding that achieves minimum latency. Design trade-offs in terms of area and throughput efficiency are explored to find the optimal architecture. The proposed Turbo decoder has been modeled using Simulink; various test cases are used to estimate the performances. The results are analyzed and achieved 50% reduction in computation time along with reduced BER (e-3).

Index Terms— 3GPP LTE, Convolutional interleaver, MAP decoder, SOVA, Turbo decoder, VLSI ASIC.

1 INTRODUCTION

9 GPP Long Term Evolution (LTE) [1], which is a set of **J**enhancements to the 3G Universal Telecommunications System (UMTS) [2], has received tremendous attention recently and is considered to be a very promising 4G wireless technology. For example, Verizon Wireless has decided to deploy LTE in their next generation 4G evolution. One of the main advantages of 3GPP LTE is high throughput. For example, it provides a peak data rate of 326.4 Mbps for a 4×4 antenna system, and 172.8 Mbps for a 2×2 antenna system for every 20 MHz of spectrum. Furthermore, LTE-Advance [3], the further evolution of LTE, promises to provide up to 1 Gbps peak data rate. The channel coding scheme for LTE is Turbo coding [4]. The Turbo decoder is typically one of the major blocks in a LTE wireless receiver. Turbo decoders suffer from high decoding latency due to the iterative decoding process, the forward backward recursion in the maximum a posteriori (MAP) decoding algorithm and the interleaving/de-interleaving between iterations. Generally, the task of an interleaver is to permute the soft values generated by the MAP decoder and write them into random or pseudo-random positions.

2 FUNDAMENTALS OF TURBO CODES

In order to explain the proposed Turbo decoder architecture, the fundamentals of Turbo codes are briefly described in this section.

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2.1 Turbo encoder structure

As shown in Fig. 1, the Turbo encoding scheme in the LTE standard is a parallel concatenated convolutional code with two 8-state constituent encoders and one convolutional interleaver [5]. The function of the convolutional interleaver is to take a block of N-bit data and produce a permutation of the input data block. From the coding theory perspective, the performance of a Turbo code depends critically on the interleaver structure [8]. The basic LTE Turbo coding rate is 1/3. It encodes an N-bit information data block in to a code word with 3N+12 data bits, where 12 tail bits are used for trellis termination. The initial value of the shift registers of the 8-state constituent encoders shall be all zeros when starting to encode the input information bits. LTE has defined 188 different block sizes.

The convolutional encoder can be represented as follows [6]:

The convolutional encoder basically multiplies the generator Polynomials by the input bit string, as follows:

•
$$A(x) = g0(x) * I(x) = a b c ... g$$

• $B(x) = g1(x) * I(x) = P Q R ... V$

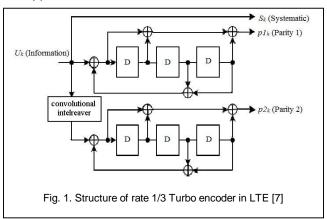
Interleaving the two outputs from the convolutional encoder yields $E(x) = aPbQcR \dots gV$, which can also be written as:

$$E(x) = (a0\ b0\ c0\ ...\ g0) + (0P0Q0R\ ...\ 0V) = A(x2) + x^*B(x2)$$

Therefore, $E(x) = A(x2) + x^*B(x2)$ and $A(x2) = g0\ (x2) + I(x2)$ and $B(x2) = g1(x2) * I(x2)$, with the following. $E(x) = g0(x2) * I(x2) + x * g1(x2) * I(x2)$

=
$$I(x2) * G(x)$$

Where $G(x) = g0(x2) + x * g1(x2)$
i.e. $G(x) = 1 + x + x2 + x4 + x5 + x6 + x7 + x11 + x12 + x13$.



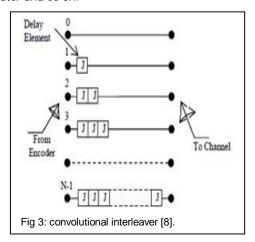
2.2 Turbo Decoder Structure

The basic structure of a Turbo decoder is functionally illustrated in Fig.2. A turbo decoder consists of two maximum a posteriori (MAP) decoders separated by an interleaver that permutes the input sequence. The decoding is an iterative process in which the so-called extrinsic information is exchanged between MAP decoders. Each Turbo iteration is divided in to two half iterations. During the first half iteration, MAP decoder 1 is enabled. It receives the soft channel information (soft value Ls for the systematic bit and soft value L_{p1} for the parity bit) and the a priori information L_{a1} from the other constituent MAP decoder through deinterleaving to generate the extrinsic information L_{e^1} at its output. Likewise, during the second half iteration, MAP decoder 2 is enabled, and it receives the soft channel information (soft value Ls for a permuted version of the systematic bit and soft value $L_{p,2}$ for the parity bit) and the a priori information L^2 from MAP decoder1 through interleaving to generate the extrinsic information L_{e^2} at its output. This iterative process repeats until the decoding has converged or the maximum number of iterations has been reached.

2.3 Convolutional Interleaver

A convolutional interleaver [8] consists of N rows of shift registers, with different delay in each row. In general, each successive row has a delay which is J symbols duration higher than the previous row as shown in Fig. 3. The code word symbol from the encoder is fed into the array of shift registers, one code symbol to each row. With each new code word symbol the commutator switches to a new register and the new code symbol is shifted out to the channel. The i-th (1 \leq i \leq N-1) shift register has a length of (i-1)J stages where J = M/N and the last row has M-1 numbers of delay elements. The convolutional deinterleaver performs the inverse operation of the interleaver and differs in structure of the arrangement of

delay elements. Zeroth row of interleaver becomes the N-1 row in the deinterleaver. 1st row of the former becomes N-2 row of later and so on.



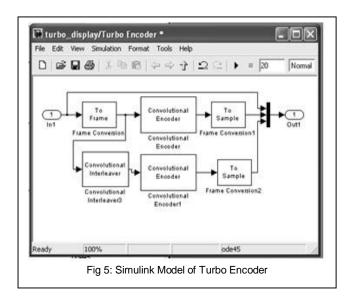
In order to verify the Verilog HDL models for the interleaver and deinterleaver the authors have developed another top level Verilog HDL model, combining interleaver and deinterleaver [8]. The scrambled code words from the output of the interleaver is applied as input to the deinterleaver block along with clock as synchronization signal. It is observed in Fig 4 that the scrambled code word is converted into its original form at the output of the deinterleaver block.

3 SOFTWARE ANALYSIS OF TURBO DECODER

Simulink model of a Turbo Encoder and Turbo Decoder are shown below.

3.1 Turbo Encoder

It consists of two convolutional encoders. The outputs of the turbo encoder are the information sequence, together with the corresponding parity sequence produced by first encoder and the parity sequence produced by the second encoder block, the input to second encoder is through interleaver, which scrambles the data bit sequence. Simulation model of Turbo encoder – decoder is shown in fig 5.

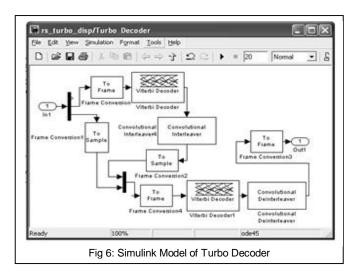


3.2 Turbo Decoder

The proposed Turbo decoder shown above in Fig 6 uses iterative decoding. The turbo code decoder is based on a modified Viterbi algorithm that incorporates reliability values to improve decoding performance. The turbo decoder consists of M-elementary decoders one for each encoder in turbo encoding part. Each elementary decoder uses the Soft Decision Viterbi Decoding to produce a soft decision for each received bit. After an iteration of the decoding process, every elementary decoder shares its soft decision output with the other (M-1) elementary decoders.

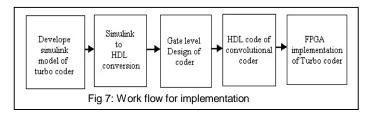
3.3 Channel

The AWGN Channel [1] block adds white Gaussian noise to a real or complex input signal. Each of the major blocks mentioned above have individual sub blocks which are configured to meet the Specifications (After scaling, keeping in mind the mathematical constraints of modeling a real time system).



4 WORK FLOW

First *Simulink* model for turbo encoder and Turbo decoder is developed on MATLAB platform. The verilog HDL coding is made on *Modelsim* and verified. Finally implemented on FPGA, Fig 7 shows the work flow for implementation.

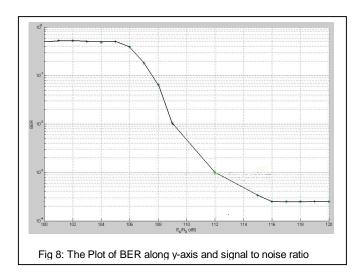


5 RESULTS AND DISCUSSIONS

With the above Turbo encoder, Turbo decoder and channel specifications, the Bit Error Rate (BER) is analyzed for 3 samples per frame; it is 0.000254 so error rate is of order (10-3). The *simulink* model for turbo encoder and turbo decoder is as shown in Fig.5 and Fig 6. Similarly with the same encoder and decoder specifications are tested for different samples per frame. Here the computation time varies as the number of samples per frame increases; these cases also provide the BER of order 10-3. The obtained results are shown in the Table 1, and the obtained graph is as shown in Fig.8.

Table1: Bits Per Frame Vs Corresponding BER

Bits Per frame	SNR(dB)	BER
3	5	0.254e-3
	10	0.214e-3
7	5	0.410e-3
	10	0.393e-3
12	5	0.324e-3
	10	0.311e-3
15	5	0.364e-3
	10	0.332e-3



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6 CONCLUSION

In this paper, presented a brief survey about turbo codes, designed a turbo encoder and turbo decoder on simulink and coded in Verilog HDL and implemented on FPGA. The codes and coding techniques are carried by Viterbi decoder. For turbo codes, turbo decoders and their decoding algorithms like log-MAP/SOVA are used. Based on the algebraic constructions, the interleaver offers capability which enables Turbo decoding by using MAP decoders working concurrently. We proposed low complexity recursive architecture for generating the convolutional interleaver addresses on the fly. The convolution interleavers are designed to operate at full speed with the MAP decoders. The proposed architecture has been scaled and can be tailored for different throughput requirements.